

# ASDF: A Compiler for Qwerty, a Basis-Oriented Quantum Programming Language

CGO '25

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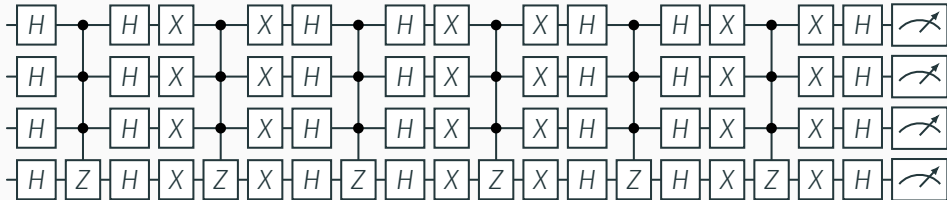
March 4th, 2025

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# Background: Quantum Computing

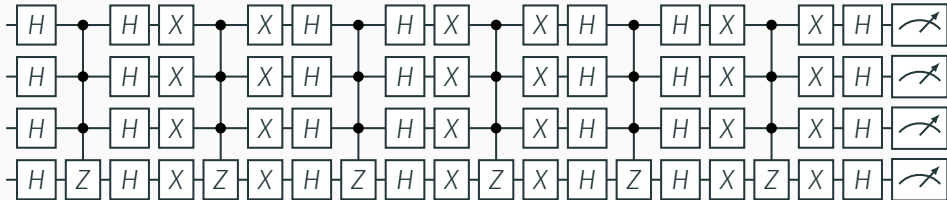
- Quantum computers promise exponential speedup for important problems (e.g., integer factoring and physics simulation)
- ...but current quantum programming languages (e.g., Q# or Qiskit) require programming in low-level quantum assembly (quantum *gates* and *circuits*)

## Background: Example Quantum Circuit



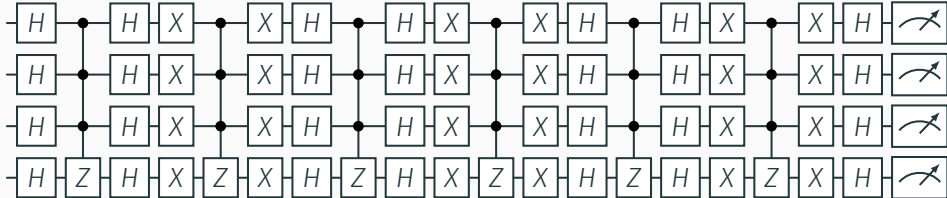
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*Unstructured search algorithm:*



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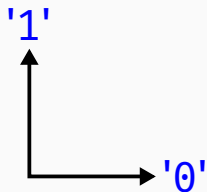
Tedious, tricky to write (like classical assembly)

## Background: The Qwerty Programming Language

- **Qwerty**: high-level quantum DSL embedded in Python
- Primitives are **basis translations** rather than quantum gates
  - Computation is a pipeline:  
 $x \mid f \mid g$  means  $g(f(x))$

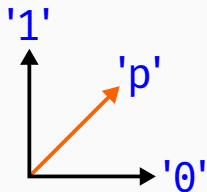
# Hello World in Qwerty

Qubit literals:



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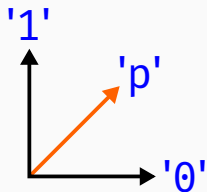
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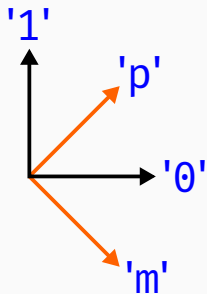


Random bit generator:

'p' | **measure**

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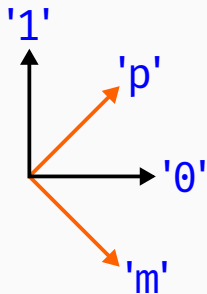
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{ 'p', 'm' }

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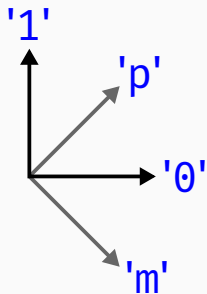
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```
'p' | {'p', 'm'} >> {'1', '0'}  
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Always measures a 1: **Basis translation**

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## Realistic Qwerty Example: Grover diffuser

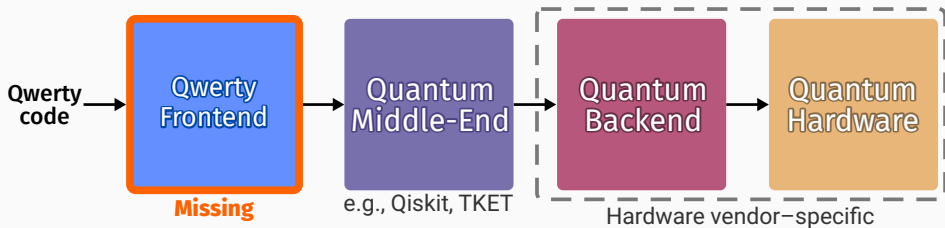
Qwerty:

```
'p'[N] >> -'p'[N]
```

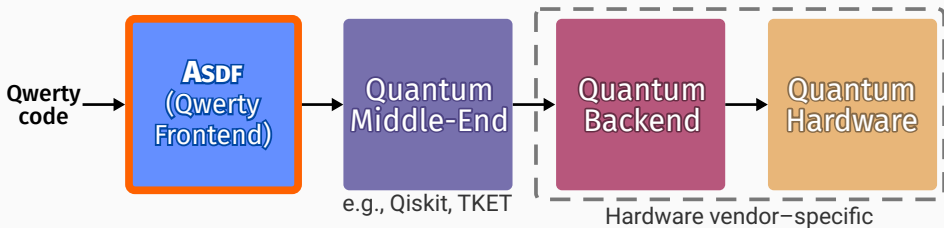
Q# (Prior work):

```
within {  
    ApplyToEachA(H, q);  
    ApplyToEachA(X, q);  
} apply {  
    Controlled Z(Most(q),  
                Tail(q));  
}
```

# Motivation: Qwerty Needs a Compiler

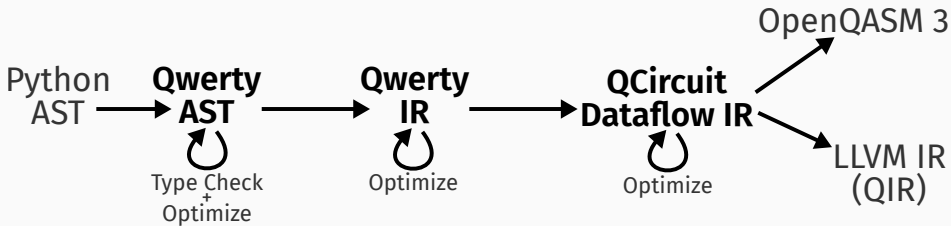


## Motivation: Qwerty Needs a Compiler



We present ASDF, the first compiler for a basis-oriented quantum programming language.

# Overview of ASDf





- ① Fast compilation of basis-oriented operations

## Challenges in Compiling Qwerty

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- ④ Integration with quantum ecosystem

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## Challenge ④ – Integration:

5. Embedded in **Python**, outputs **industry-standard IRs**



# Span Equivalence Checking

- Core Qwerty primitive: **basis translation**  $b_1 \gg b_2$ , where  $b_1$  and  $b_2$  are bases
- Qwerty type checking requires that  $\text{span}(b_1) = \text{span}(b_2)$ 
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✓ Spans are equivalent

ASDF checks span equivalence in  $O(n^2 \log n)$  time instead of exponential time

- Qwerty IR is the quantum MLIR dialect with the highest level of abstraction
- For example, `'p'[3] >> -'p'[3]` becomes the following IR:

```
%12 = arith.constant 3.14159
```

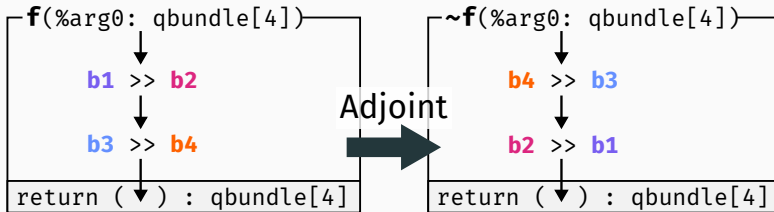
```
%13 = qwerty.qbtrans %8 by {"ppp"} >> {exp(i*%12)*"ppp"}
```





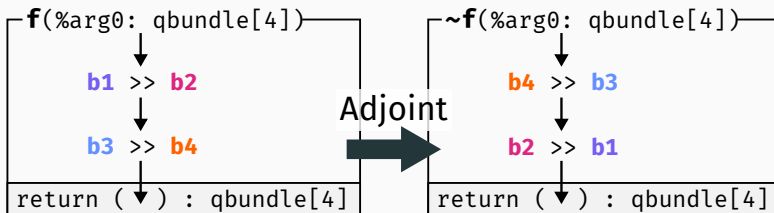
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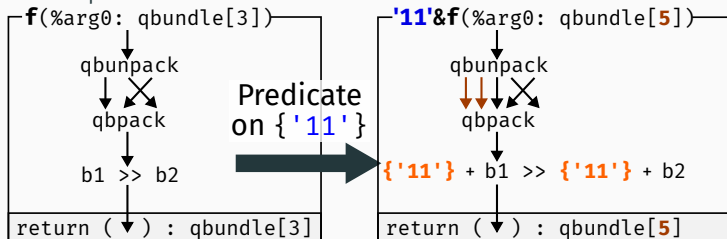


- Novel `Adjointable` op interface in MLIR

# Predicating Basic Blocks

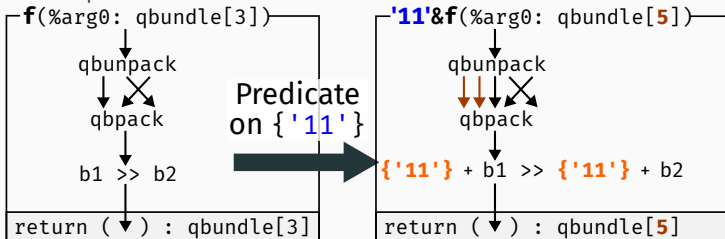
- Qwerty syntax for *predicating* a function  $f$  with basis  $b$ :  
     $b \ \& \ f$
- $b \ \& \ f$  will run only in the proper subspace  $b$

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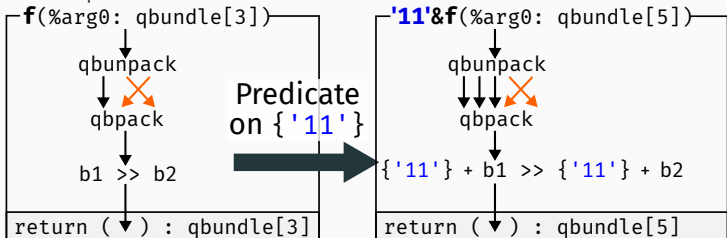
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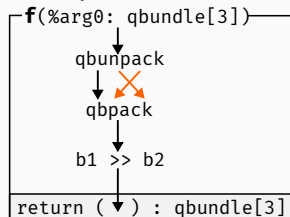
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## Qwerty Code

```
@qpu
def f(q: qubit[3]) -> qubit[3]:
  q1, q2, q3 = q
  return q1+q3+q2 | b1 >> b2
```

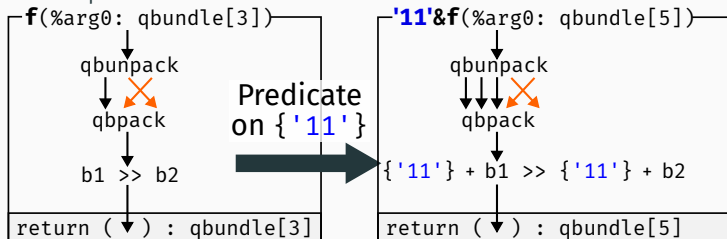


## Qwerty IR



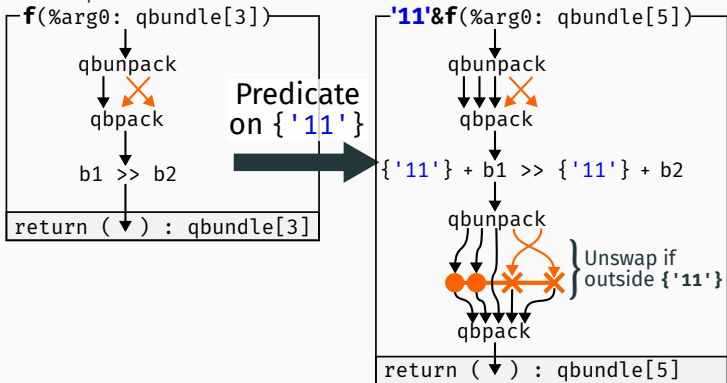
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# Basis Translation Synthesis: Example 1

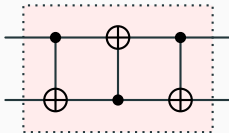
$\{ '01', '10' \} \gg \{ '10', '01' \}$



Permutation

$|00\rangle \mapsto |00\rangle$      $|01\rangle \mapsto |10\rangle$

$|10\rangle \mapsto |01\rangle$      $|11\rangle \mapsto |11\rangle$



# Basis Translation Synthesis: Example 1

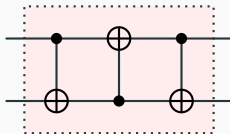
`{'01', '10'} >> {'10', '01'}`



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Permutation synthesis uses Tweedledum library from EPFL

## Basis Translation Synthesis: Example 2

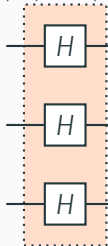
'p'[3] >> -'p'[3]



Standardize

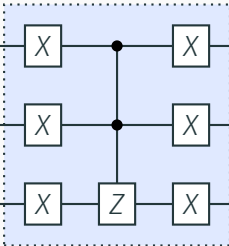
$|+\rangle \mapsto |0\rangle$

$|-\rangle \mapsto |1\rangle$



Vector phase

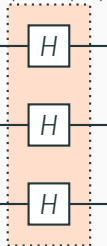
$|000\rangle \mapsto -|000\rangle$



Destandardize

$|0\rangle \mapsto |+\rangle$

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## Basis Translation Synthesis: Example 2

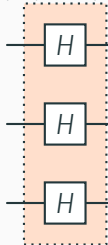
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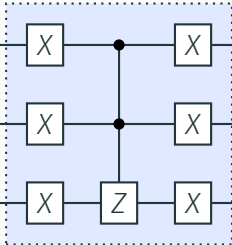
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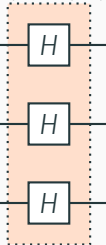
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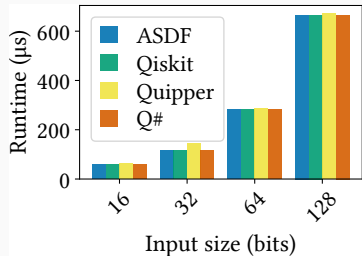


ASDF is the first compiler capable of synthesizing quantum circuits from basis translations

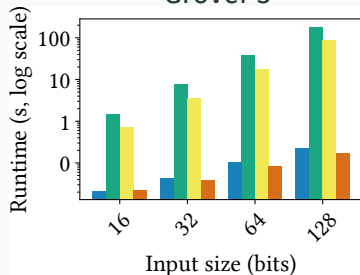
How do ASDF-synthesized circuits compare to handwritten circuits?

# Evaluation: Fault-Tolerant Runtime

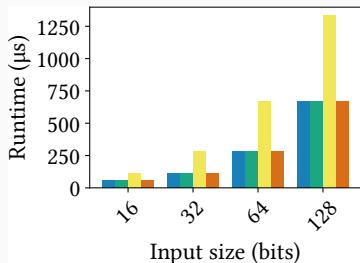
## Bernstein–Vazirani



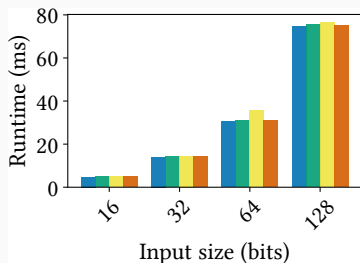
## Grover's



## Simon's

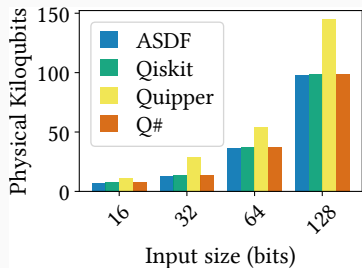


## Period finding

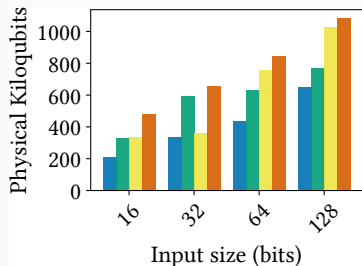


# Evaluation: Fault-Tolerant Physical Qubits

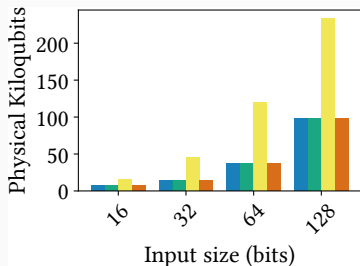
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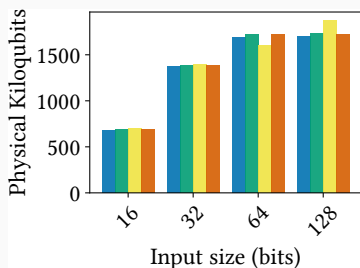
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## Period finding



Overall, AsDF keeps pace with handwritten circuits compiled with gate-oriented compilers.



## Conclusion

In this talk, I presented **ASDF**, a compiler that leverages novel basis-oriented compilation techniques to enable Qwerty's high-level quantum programming paradigm with minimal overhead.

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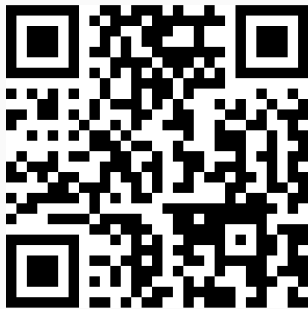
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Qwerty tech report:



[arXiv:2404.12603](https://arxiv.org/abs/2404.12603)

Source code:



[github.com/gt-tinker/qwerty](https://github.com/gt-tinker/qwerty)

## Backup Slides

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## Full Bernstein–Vazirani Example Program

```
1 from qwerty import *
2
3 def bv(secret_string):
4     @classical[[N]](secret_string)
5     def f(secret_string: bit[N], x: bit[N]) -> bit:
6         return (secret_string & x).xor_reduce()
7
8     @qpu[[N]](f)
9     def kernel(f: cfunc[N,1]) -> bit[N]:
10        return 'p'[N] | f.sign \
11                | pm[N] >> std[N] \
12                | measure[N]
13
14    return kernel()
15
16 secret_string = bit.from_str('1101')
17 print(bv(secret_string))
```

## Predication Example

Imagine '0' & ({'0', '1'} >> {'p', 'm'}).

This performs the following:

'00' ↦ '0p'

'01' ↦ '0m'

'10' ↦ '10'

'11' ↦ '11'

## QCirc Dialect Example

```
1 %q = qcirc.H %0
2 %q_0 = qcirc.H %1
3 %q_1 = qcirc.H %2
4 %q_2 = qcirc.X %q_1
5 %q_3 = qcirc.X %q_0
6 %q_4 = qcirc.X %q
7 %ctrlq:2, %q_5 = qcirc.Z [%q_4, %q_3] %q_2
8 %q_6 = qcirc.X %q_5
9 %q_7 = qcirc.X %ctrlq#1
10 %q_8 = qcirc.X %ctrlq#0
11 %q_9 = qcirc.H %q_8
12 %q_10 = qcirc.H %q_7
13 %q_11 = qcirc.H %q_6
```

Inspired by QIRO and QSSA

# Qubit Index Analysis

